

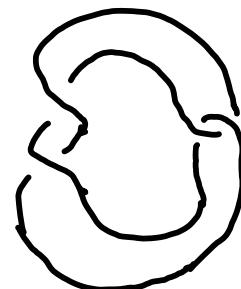
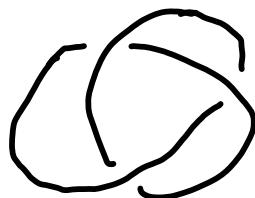
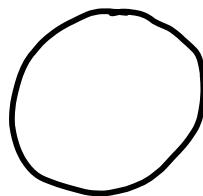
Knot Theory and its Combinatorial Invariants

Shelly Harvey
Rice University

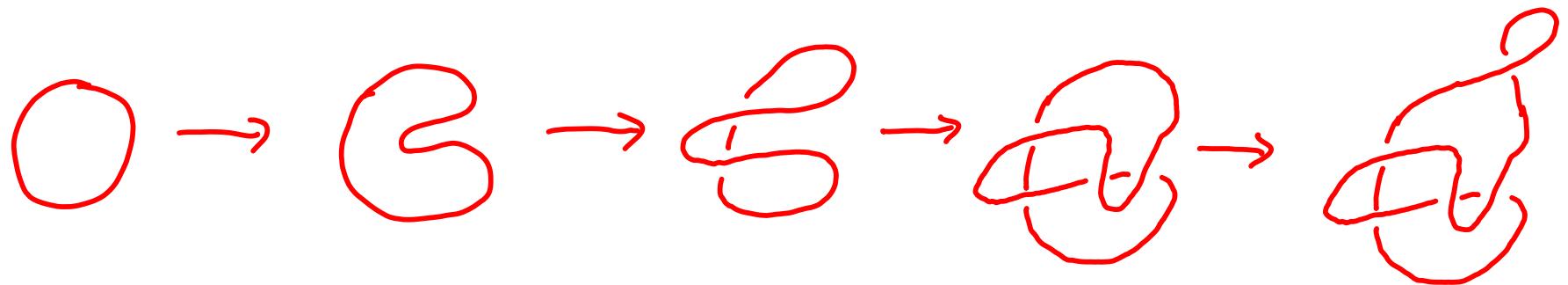
Def: A knot is a smooth embedding

$$f: S^1 \rightarrow \mathbb{R}^3.$$

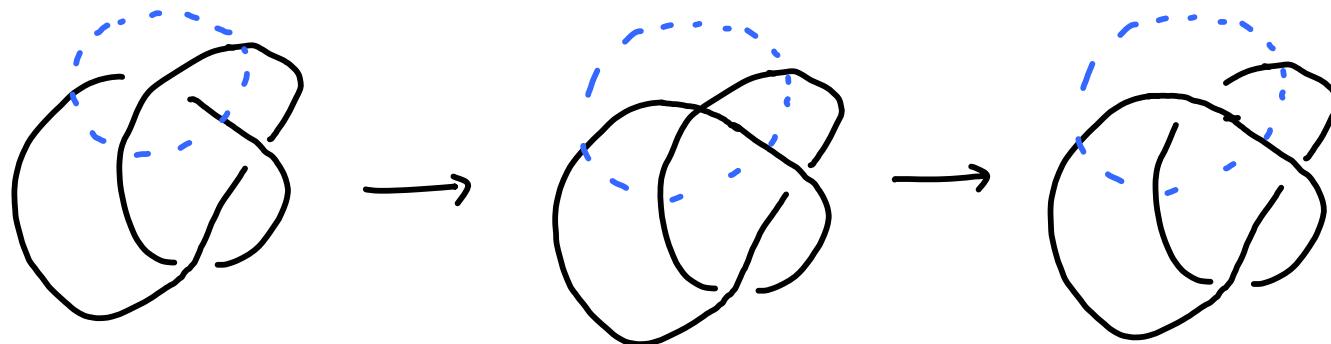
i.e. take a rope, tie it up and attach
the ends.



We say that two knots are equivalent if you can deform one into the other without allowing the curve to pass through itself.

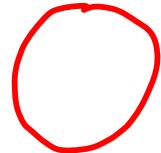


all equivalent knots

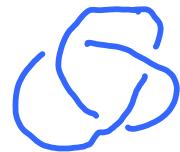


NOT ALLOWED

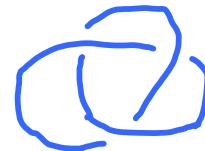
Examples of Distinct Knots:



unknot



left-handed
trefoil



right-handed
trefoil

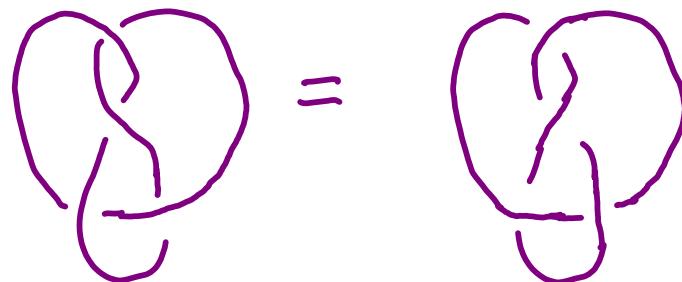
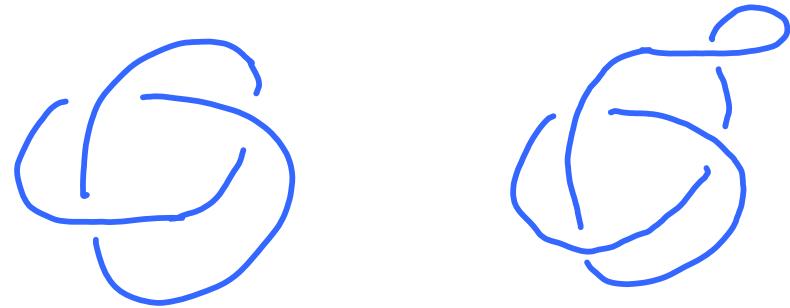


figure-eight

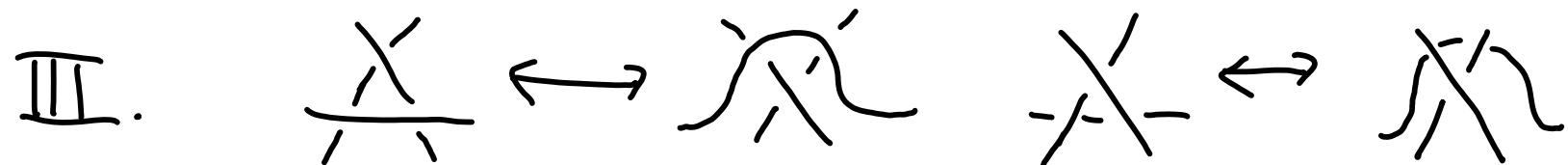
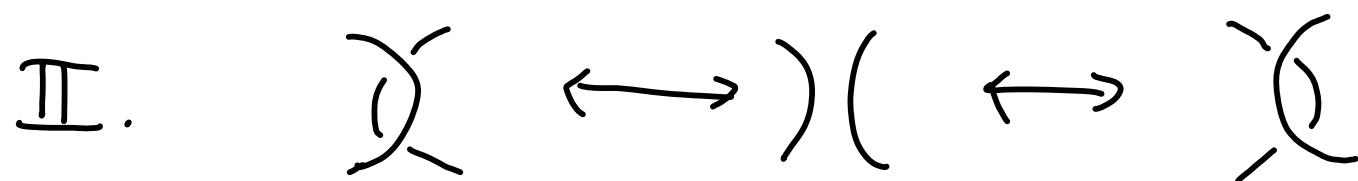
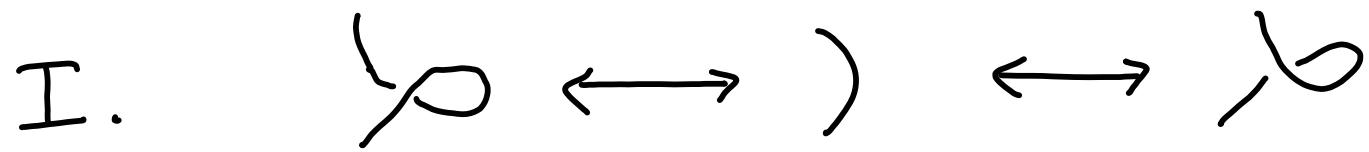
Instead of thinkings of knots as subsets of \mathbb{R}^3 , we project them to plane and remember the crossing information.

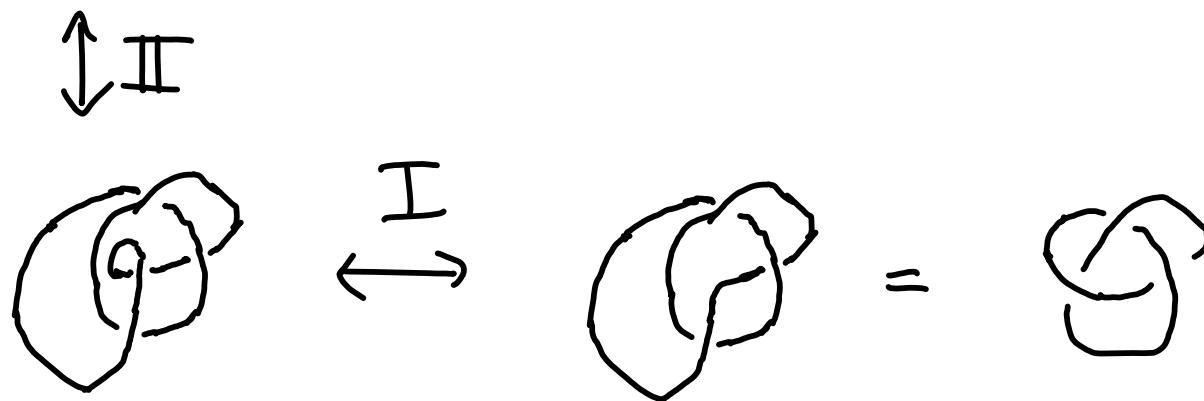
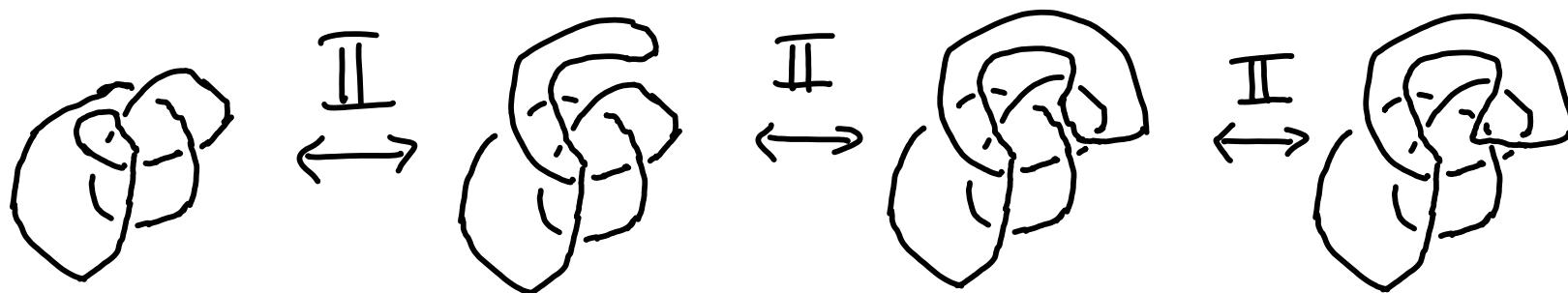
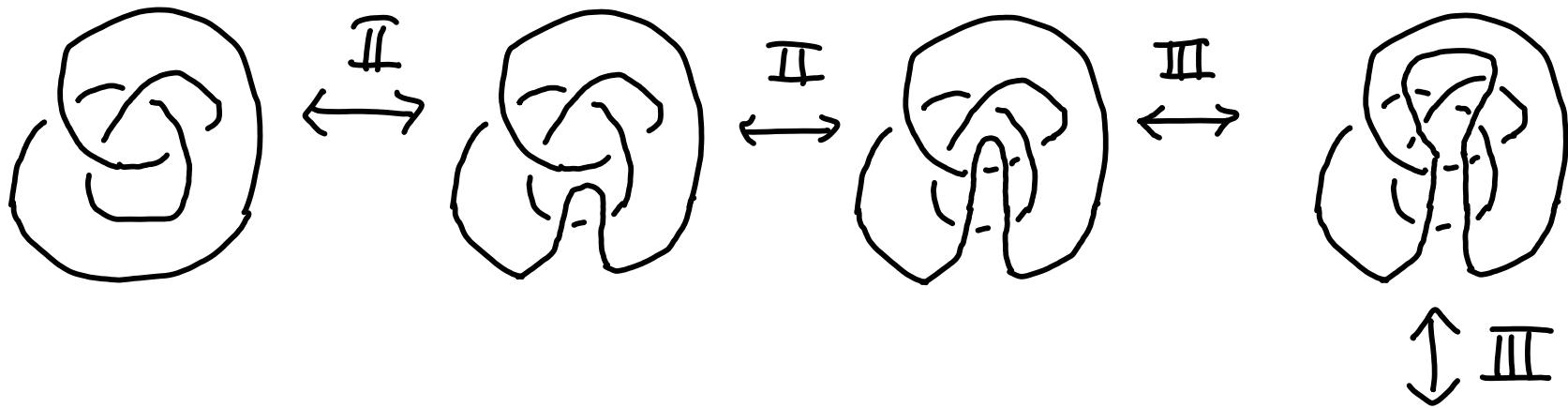
This is called a diagram of the knot.



Different diagram for the same knot

Theorem (Alexander, Briggs): Two diagrams are associated to the same knot \Leftrightarrow they are related by a sequence of the Reidemeister moves (plus planar isotopies).





[many diagrams
for same knot!]

Thus to define a knot invariant, it suffices to consider invariants of diagrams that are preserved under Reidemeister moves:

$$I: \{ \text{diagrams} \} \longrightarrow A = \left\{ \begin{array}{l} \text{rings,} \\ \text{algebra, etc.} \end{array} \right\}$$

$I(D_1) = I(D_2)$ if D_1 is related to D_2 by Reidemeister moves I, II, III.

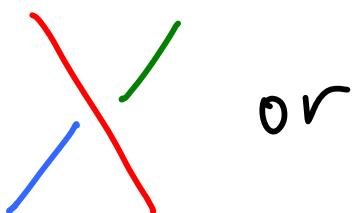
The simplest invariant is 3-colorability:

$$G: \{\text{Diagrams}\} \rightarrow \{0, 1\}.$$

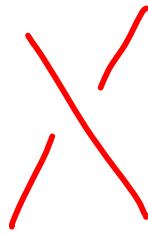
$c(D) = 1$ if D is 3-colorable, i.e.

the arcs in the projection can be colored with $\{R, G, B\}$ s.t. every colored is used and the crossings look like :

(all 3 colors
appear)



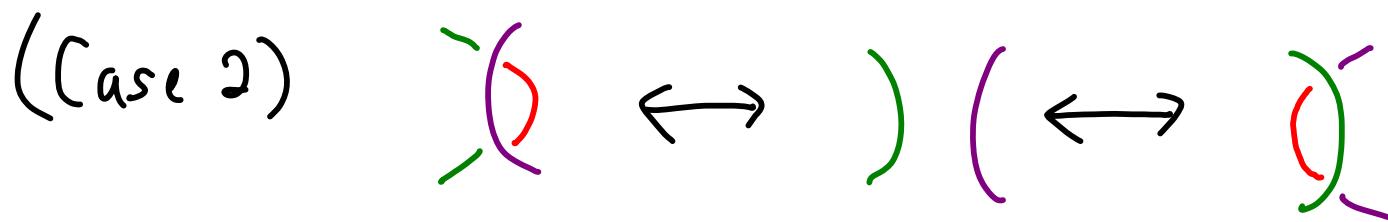
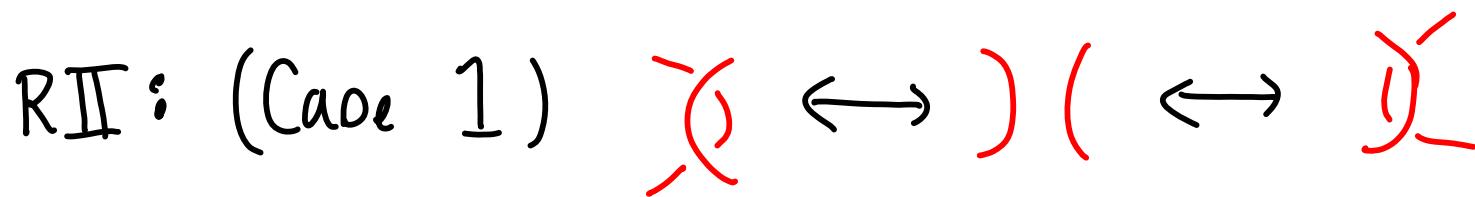
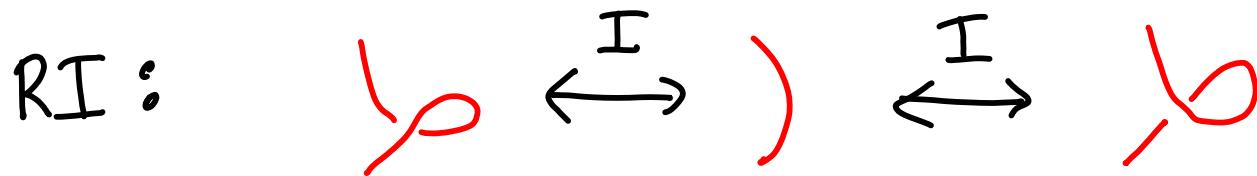
or



(all one color)

Otherwise $c(D) = 0$.

Can check that if D_1 and D_2 differ by Reidemeister moves, then $c(D_1) = c(D_2)$.

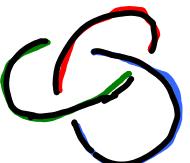


R~~II~~ III is similar

Therefore 3-colorability is a property of
the knot, not just its diagram.

Def $c(K) := c(D)$ for some diagram
 ↑ ↑
 knot D of K.
 $K: S^1 \rightarrow \mathbb{R}^3$ D = a diagram or
 "picture" of K

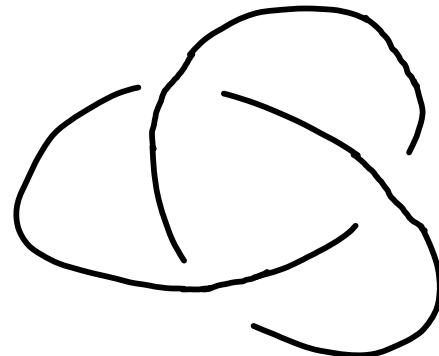
Ex:  is not 3-colorable $\Rightarrow C(O) = \emptyset$

 is 3-colorable $\Rightarrow C(S) = 1$

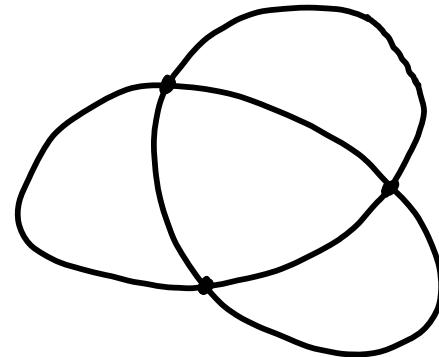
$\therefore O \neq S$ (as knots)

Let K be a knot and D be a diagram for K . If we forget about whether crossings are over/under, we get a planar graph :

crossings \leftrightarrow vertices

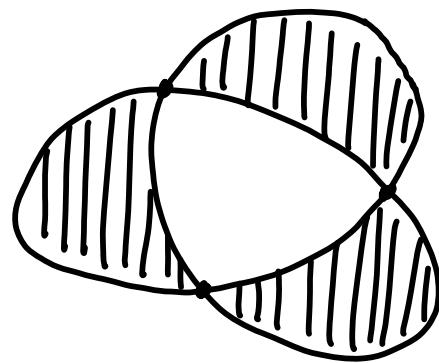


D = diagram

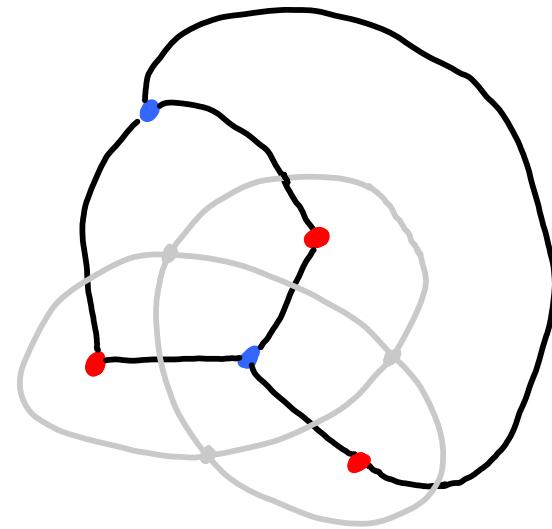


P_D = projection,
4-regular graph

The dual of P_D is a planar bipartite graph, each of whose regions has 4 edges.



Checkerboard
coloring of P_D

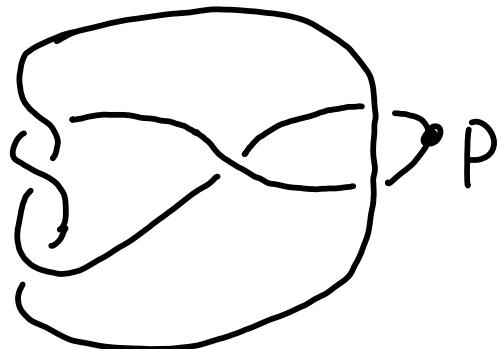


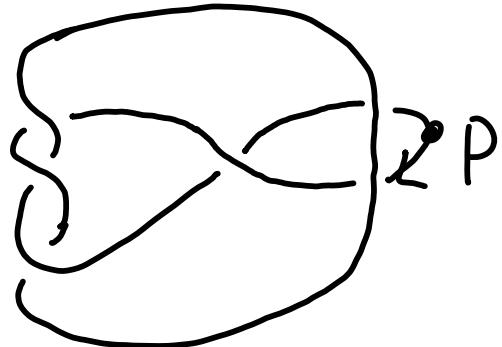
Red vertices \leftrightarrow

Blue vertices \leftrightarrow

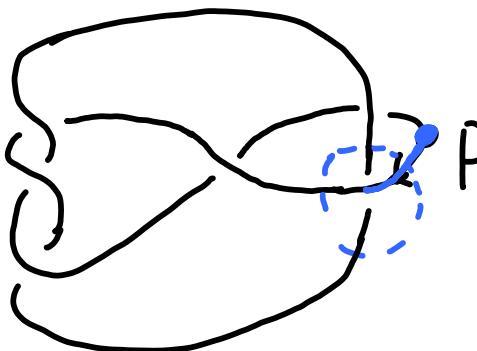
Fact: every knot can be made into the unknot by changing a subset of crossings in a diagram

Proof: Let D be a diagram for K .

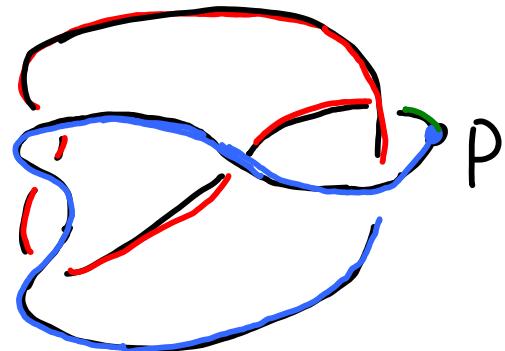




Start at P , traverse the knot.
Each time you hit a crossing,
if it is the first time you
hit a crossing, make sure you are crossing
over the other arc (if not change crossing).

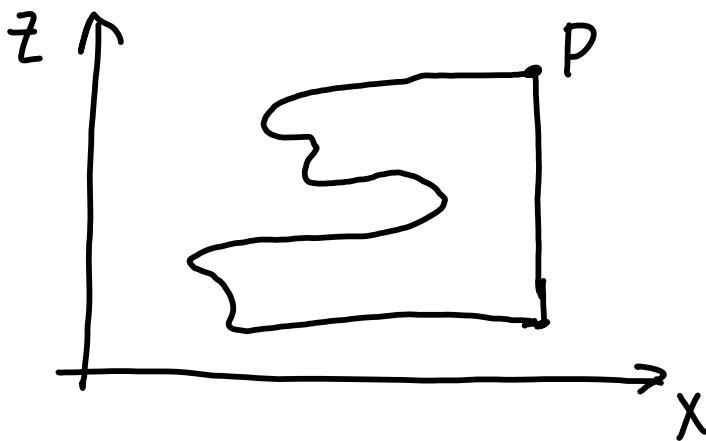


$\rightarrow \dots \rightarrow$



Then the z-value of the points on knot decreases as you traverse the knot except near p where there is a nearly vertical arc.

From side :



This is always a diagram for the unknot (O).

图

Unknotting Number

Def: $u(D) = \min \# \text{ of crossings of } D$
needed to change D to a
diagram of unknot.

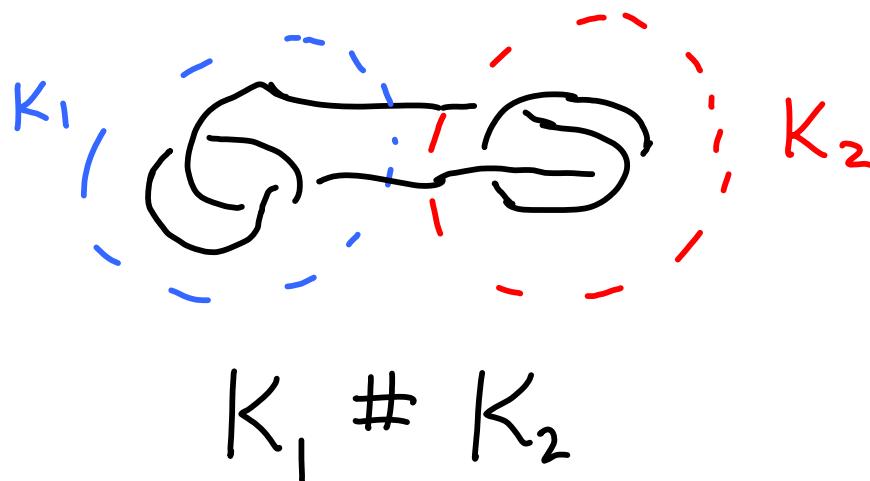
$$u(K) = \min \{ u(D) \mid D \text{ is a diagram for } K \}$$

$u(K)$ is called the unknotting number of K .

Unless $u(K)=1$ or $u(K)=0$, it is very
difficult to determine $u(K)$!

Open

Conjecture: $u(K_1 \# K_2) = u(K_1) + u(K_2)$

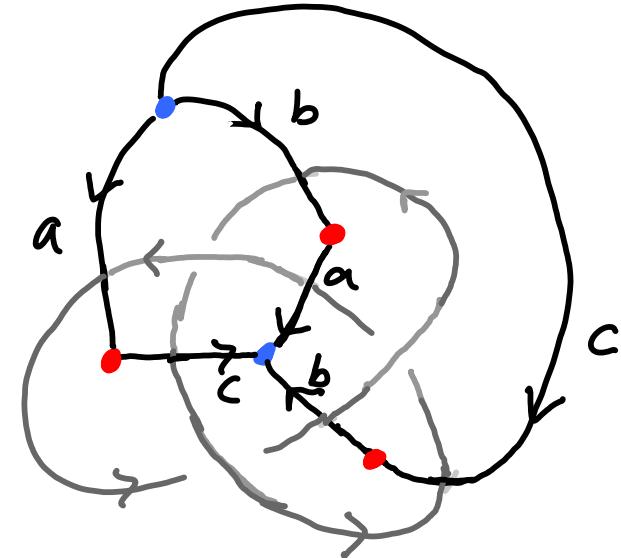
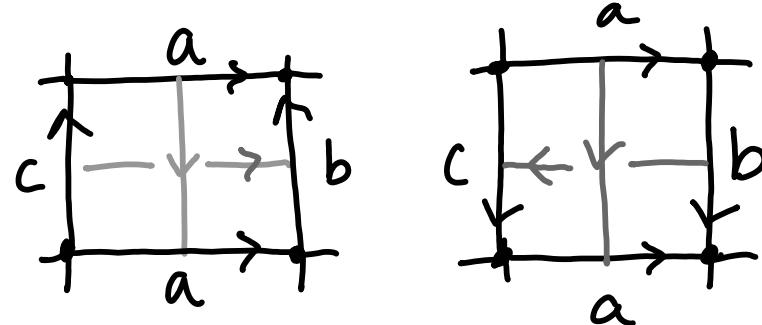


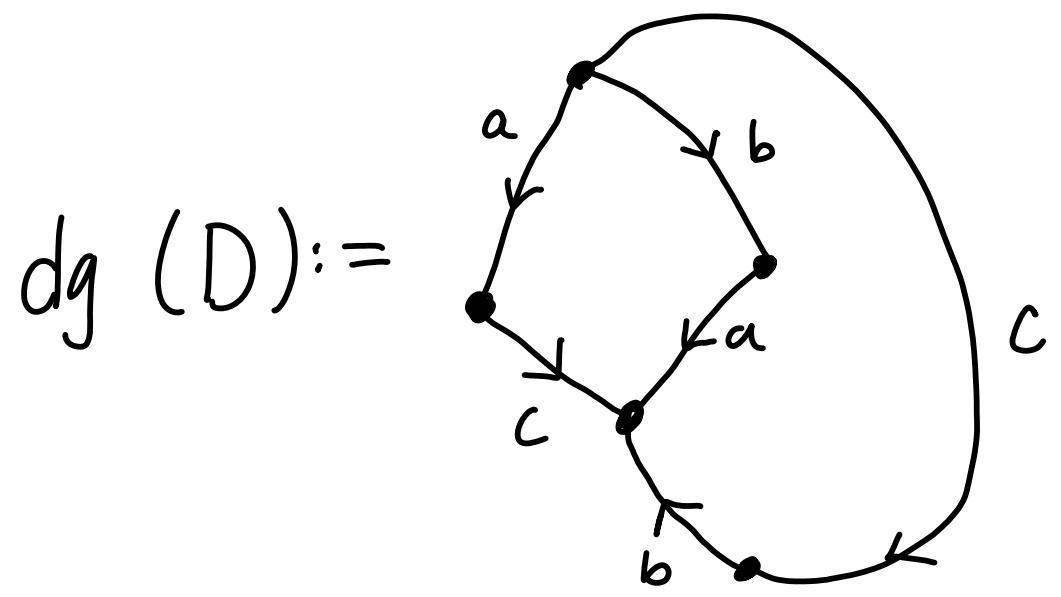
Easy to see that $u(K_1 \# K_2) \leq u(K_1) + u(K_2)$.

Since any knot can be changed to the unknot by changing crossings, to get knotting information, we need to remember crossing information.

To do this, we consider (oriented) colored graphs :

- orient knot
- orient and color edges w/ a, b, \dots s.t.



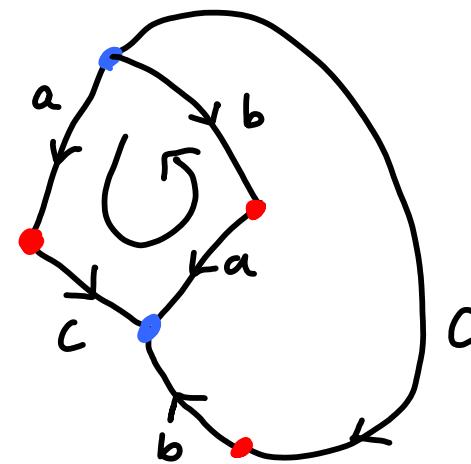


Note: Can recover the knot from the colored graph (orientations on edges gives an orientation on knot).

Define $G(D) := \langle a, b, c, \dots \mid \text{relations coming from}$
 compact regions in $dg(D)$ \rangle
 (finitely presented group)

Ex:

$$dg(D) :=$$



$$G(D) = \langle a, b, c \mid aca^{-1}b^{-1}, ab^{-1}c^{-1}b \rangle$$

It is easy to show that if D_1 and D_2 are related by Reidemeister moves then $G(D_1) \cong G(D_2)$.
as groups

Hence $G(K) := G(D)$, for any diagram D , is a knot invariant, called the knot group.

In fact, $G(K) \cong \pi_1(\overset{\uparrow}{\mathbb{R}^3 - K})$
independent of diagram.

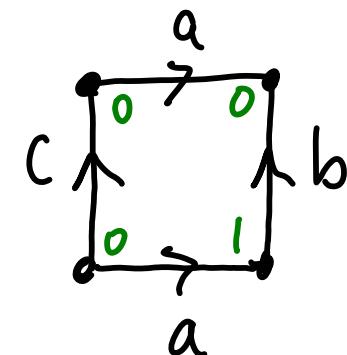
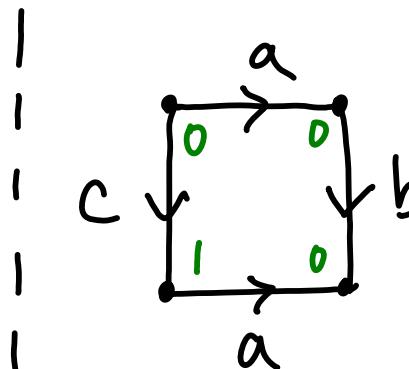
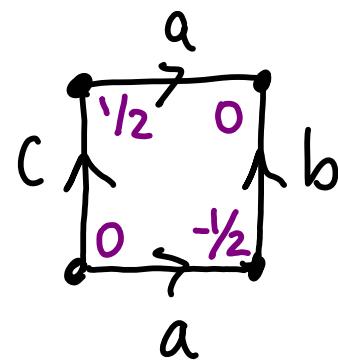
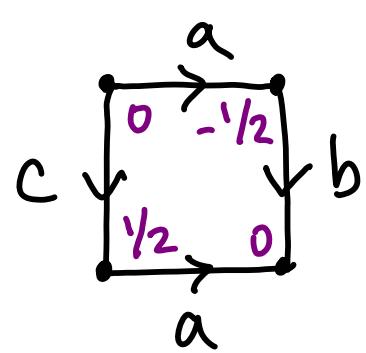
Theorem (Waldhausen, Gordon-Luecke): If there is an isomorphism $\varphi: G(K_1) \rightarrow G(K_2)$ that preserves the "peripheral structure" of K_1 and K_2 then $\underline{K_1 = K_2}$ (up to mirror image).

Thus the group of a knot is a complete invariant of knot (up to mirrors).

However, it is difficult to distinguish finitely presented groups!

Alexander Polynomial

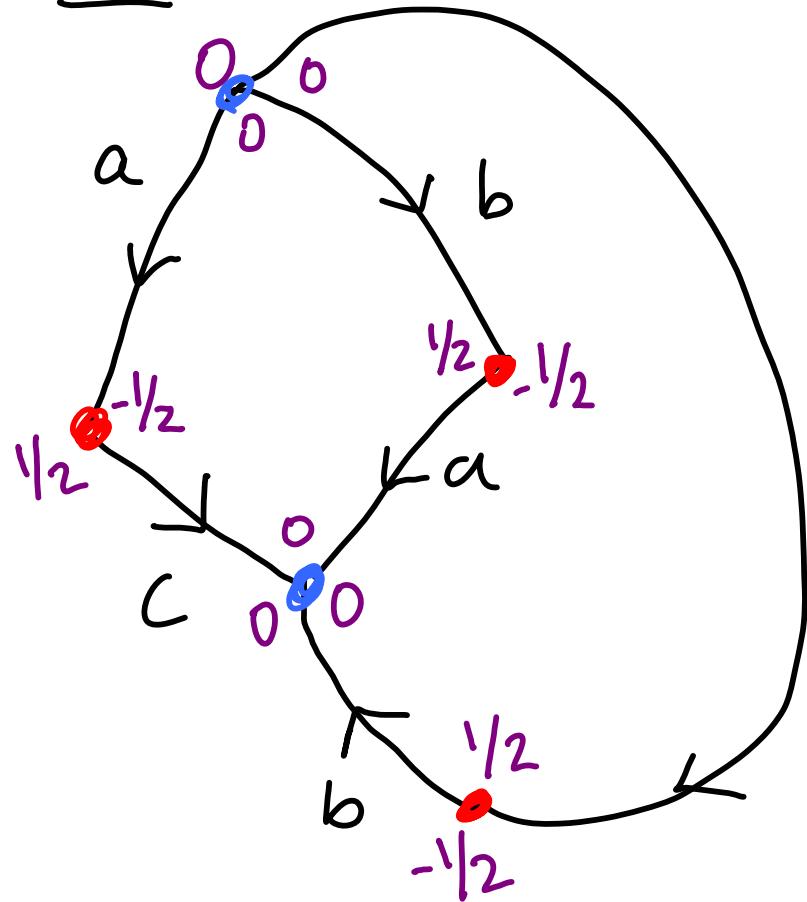
To each corner \downarrow_v in a region of $dg(D)$ assign values $A(v)$ and $B(v)$ as follows:



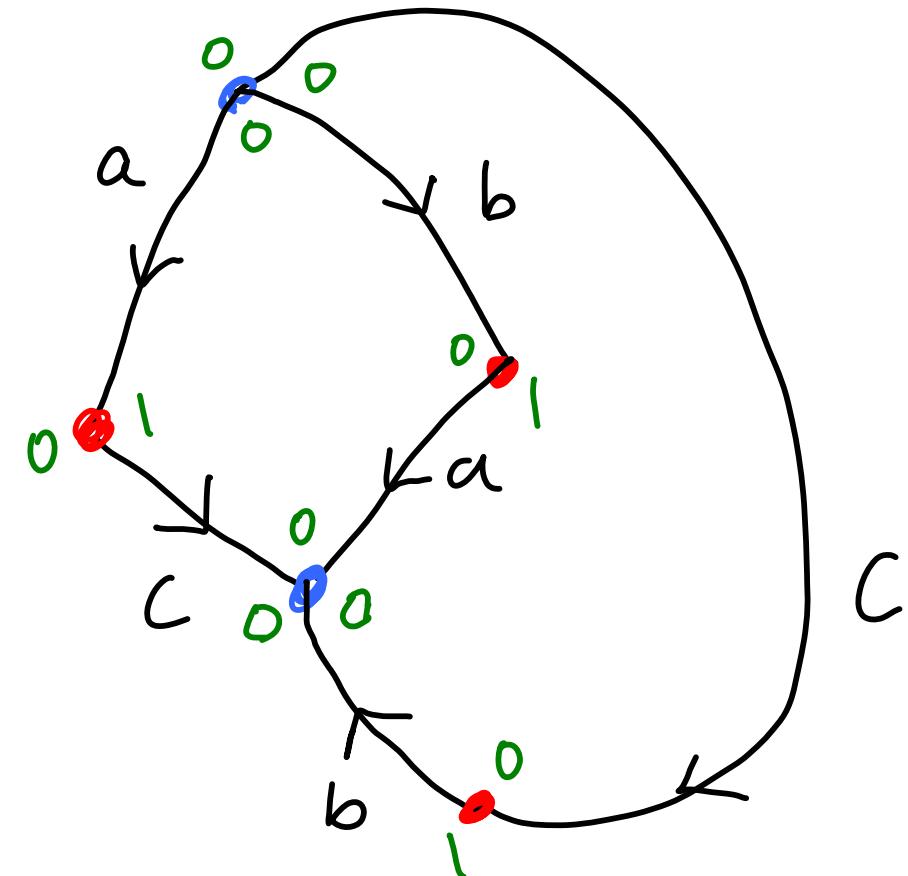
A-values

B-values

Ex:

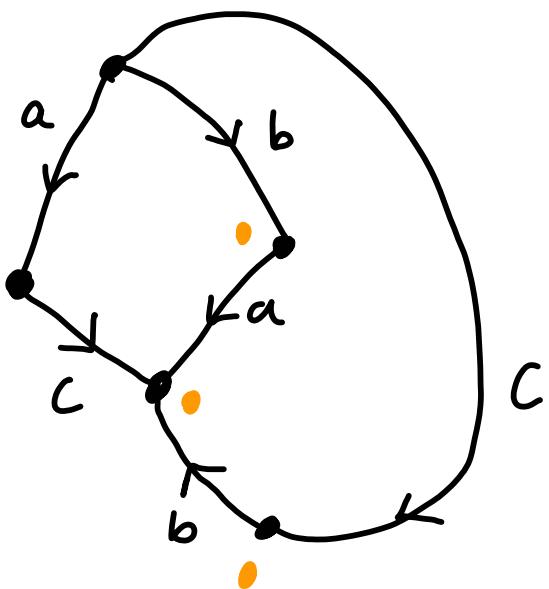


A-values



B-values

A state^(Kauffman) of $dg(D)$ is a choice of one vertex per region s.t. no vertex is used twice.



the orange dots
constitute a
single state.

Let s be a state. Then

$$A(s) := \sum_{v \in \{\text{corners in } s\}} A(v) \quad B(s) := \sum_{v \in \{\text{corners in } s\}} B(v)$$

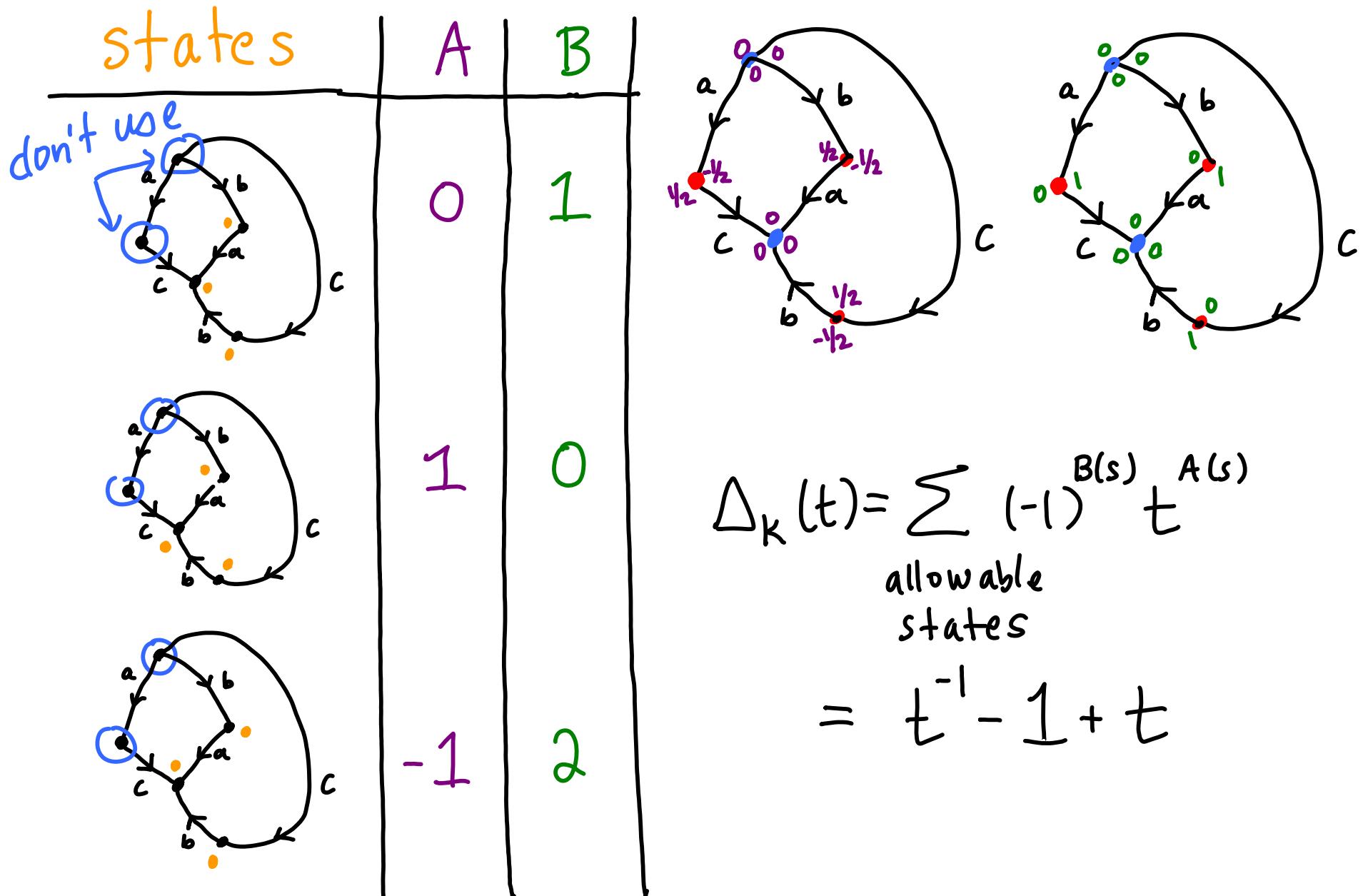
Let v_1 and v_2 be vertices of $dg(G)$ that share an edge $v_1 \xleftarrow{} v_2$.

Def: $\Delta_D(t) = \sum_{\substack{\{\text{states that do} \\ \text{not contain } v_1 \text{ or } v_2\}}} (-1)^{B(s)} t^{A(s)}$

If D_1 and D_2 are related by Reidemeister moves then $\Delta_{D_1}(t) = \Delta_{D_2}(t)$. Define $\Delta_K(t) := \Delta_D(t)$ for any diagram D for K .

$\Delta_K(t)$ is called the Alexander polynomial of K and was originally defined* by Alexander in 1928.

* Alexander gave an equivalent but different definition.



Categorifications

In 2002, P. Ozsváth - Z. Szabó and J. Rasmussen defined a bi-graded homology theory that is a categorification of the Alexander polynomial.

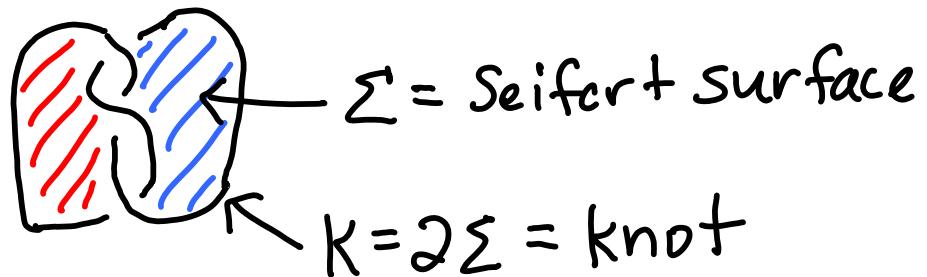
(i.e. the Euler characteristic is $\Delta_K(t)$).

knot $K \rightarrow \widehat{HFK}_j(K, i)$ abelian group $i, j \in \mathbb{Z}$
(Knot Floer homology of K)

$$\chi(\widehat{HFK}_j(K, i)) := \sum_{i,j} (-1)^j \text{rank } \widehat{HFK}_j(K, i) = \Delta_K(t)$$

Let K be a knot. A Seifert surface for K is an oriented surface embedded in $\mathbb{R}^3 - K$ whose boundary is K .

$$g(K) := \min \left\{ \frac{1 - \chi(\Sigma)}{2} \mid \begin{array}{l} \Sigma \text{ is a Seifert surface} \\ \text{for } K \end{array} \right\}$$



$$g(K) = 0 \iff K = \text{unknot}$$

Thm (Ozsváth-Szabó):

$$g(K) = \max i \text{ s.t. } \widehat{HFK}_*(K, i) \neq 0.$$

Thm (Ozsváth-Szabó, Ghinni-Ni 06):

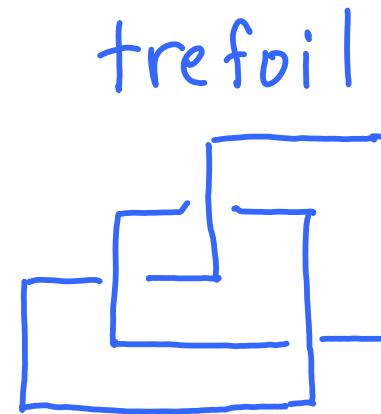
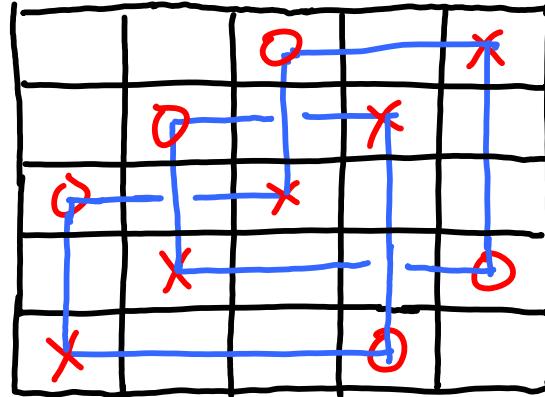
$$K \text{ is fibered} \iff \text{for } \max i \text{ s.t. } \widehat{HFK}_*(K, i) \neq 0 \\ \text{rank}(\widehat{HFK}_*(K, i)) = 1$$

The chain groups are generated by Kauffman states.

However, the boundary map is defined by counting
pseudo holomorphic curves so we cannot
determine the groups for an arbitrary knot.

Grid Diagrams

		O	X
	O		X
O		X	
	X		O
X			O



Grid diagram

One X and O
per row/column



knot

Connect X to O in columns
 O to X in rows

vertical strands cross over
horizontal strands

Define chain complex \tilde{C}_K over $\mathbb{Z}/2$.

$n!$ generators : Matchings between horizontal and vertical gridlines

n			O	X
:			O	X
:			O	X
2			X	O
1			X	

one • on each horizontal and vertical gridline (except top and right).

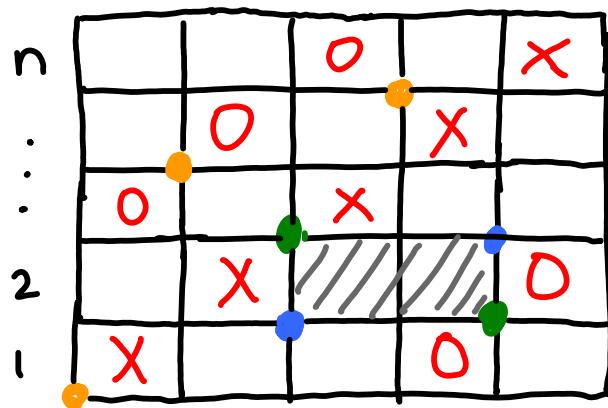
We view grid as living on a torus by identifying left/right and top/bottom.

Generators $\longleftrightarrow S_n$

$$\partial: \widetilde{C}\bar{K} \rightarrow \widetilde{C}\bar{K}$$

$$\partial X = Y + \dots$$

Counts all ways to switch SW-NE corners
of an empty rectangle to a NW-SE corner.



$$X = \bullet$$

$$Y = \bullet$$

- all points where x, y agree

M (j -grading) :

$M(x) - M(y) = (1 - \# \text{ of } 0\text{'s})$ in rectangle
with x and y corners

A (i -grading) : $(\# \text{ of } X\text{'s} - \# \text{ of } 0\text{'s})$ in
a rectangle w/ x and y
corners.

λ preserves i -grading and lowers
 j -grading by 1.

Theorem (Manolescu-Ozsváth-Sarkar):

If G is a grid diagram for K ,

$$H_*(\widetilde{C\bar{K}}(G)) \cong \widehat{HFK}(K) \otimes V^{\otimes n-1}$$

where $V = \mathbb{Z}/2_{(0,0)} \oplus \mathbb{Z}/2_{(-1,-1)}$

Hence $\chi(H_*(\widetilde{C\bar{K}}(G))) = \Delta_K(t) (1-t)^{n-1}$.

This gives the simplest algorithm to compute knot genus!

If $P_K(t)$ is any knot invariant,
a categorification of $P_K(t)$ is a bi-graded
chain complex $(C_{i,j}(K), \partial)$ s.t.

$$\begin{aligned}
\chi(C_{i,j}(K), \partial) &= \sum_{i,j} (-1)^j \text{rank } C_{i,j}(K) \cdot t^i \\
&= \sum_{i,j} (-1)^j \text{rank } H_{i,j}(K) \cdot t^i \\
&\quad \uparrow \text{homology of } (C_{i,j}(K), \partial) \\
&= P_K(t).
\end{aligned}$$

Other Categorifications

Knot Polynomial	Categorification
Alexander poly	knot Floer homology (Heegaard Floer) ~2002
Jones poly.	Khovanov homology (Khovanov) ~1999
HOMFLY poly.	Khovanov-Rosicky hom. ~2004

Categorification of graph polynomials

Polynomial	Categorification
chromatic polynomial	Helme-Guizon + Rong ~2005
Tutte polynomial	Jasso-Hernandez + Rong ~2006
Yamada polynomial	Vershinov + Vesnin ~2006

Thank you!